# X1E Architecture An Overview of Phoenix

#### NCCS USERS MEETING



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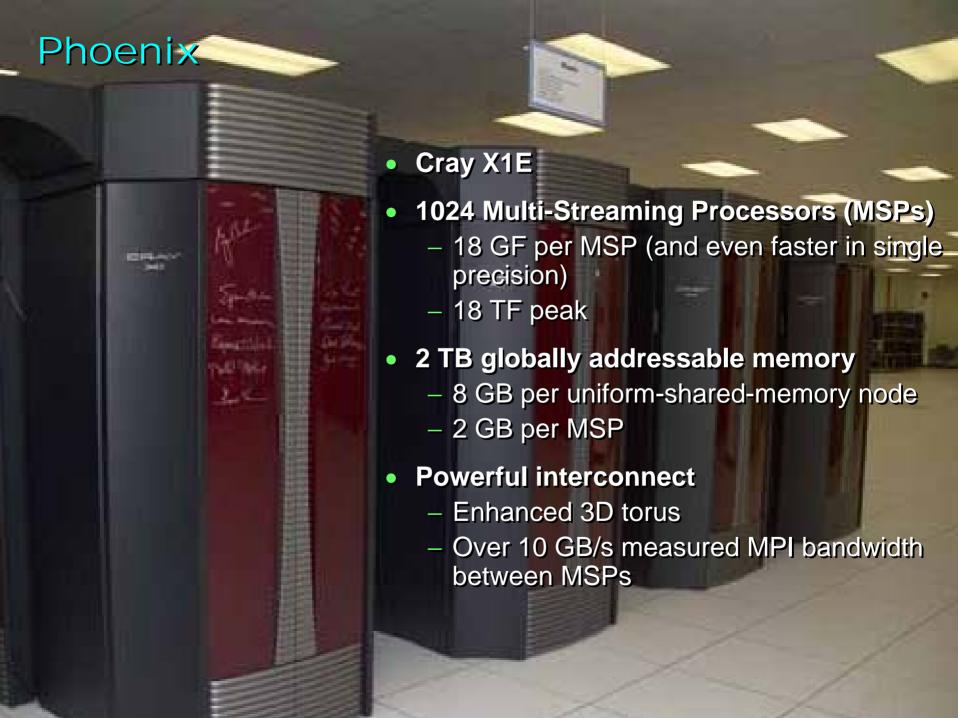
# Acknowledgements

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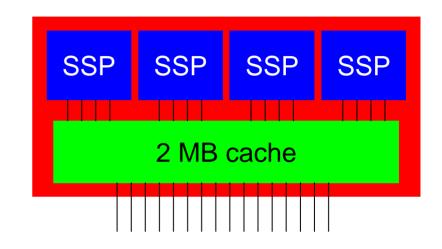
### Outline

- X1E architecture
- Using Phoenix
- Tuning is required



#### MSP versus SSP

- Each MSP has four Single-Streaming Processors (SSPs)
- Which is "the" processor?
- MSP?
  - 8-pipe vector processor
  - One MPI task
  - Automatic multi-streaming by compiler
  - 2 MB shared cache
  - Most-common mode for real applications



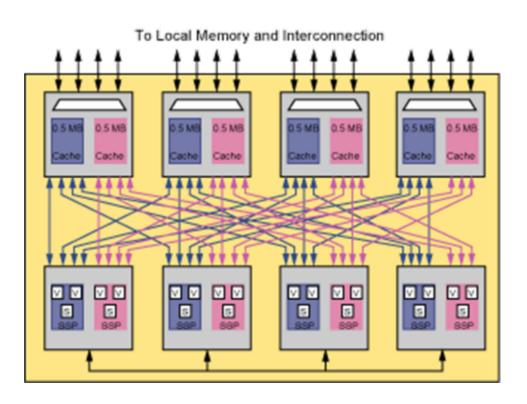
#### SSP?

- 2-pipe vector processor
- Can be an independent MPI process (or OpenMP thread)



# MCMs, nodes, and modules! Oh my!

- Node = 4 MSPs with uniform shared memory
- MCM = Multi-Chip Module = 2 MSPs in different nodes
- Module = 1 physical board = 2 nodes
  - Nodes are interleaved in hardware
  - Separate memory (still globally addressable)
  - Shared memory bandwidth
  - Shared interconnect bandwidth





#### Interconnect

#### Enhanced 3D torus

Fully connected in one dimension

#### High bandwidth

- 10.8 GB/s measured MPI point-to-point
- Takes four MSPs to saturate module bandwidth

#### Globally addressable memory

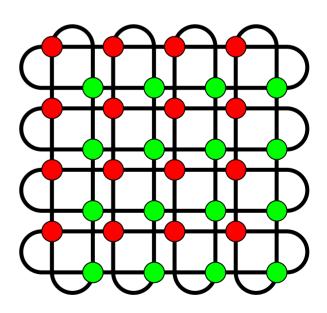
Load/store memory on any node

#### Remote address translation

- On memory's node, not at processor
- Avoids TLB misses
- Requires contiguous processors (default)
- This is why jobs migrate

#### Cache coherent

Only cache local memory



# Many levels of parallelism

- Vectorization within SSP
- Multistreaming within MSP
- OpenMP within node (not recommended)
- Between nodes (or processors)
  - MPI-1 two-sided message passing
  - MPI-2 one-sided communication
  - SHMEM one-sided communication
  - Co-Array Fortran remote memory
  - Direct load/store using pointers

# X1E strengths

#### Fast vector processors

- 18 GF double-precision peak (MSP), 15.3 GF DGEMM
- Double-rate single precision

### High memory bandwidth (local and remote)

- Stream triad of 32.7 GB/s on one MSP
- Stream triad of 15.4 GB/s/MSP fully loaded
- 10.8 GB/s MPI ping pong
- 108 GB/s Parallel Transpose
- Good at stride-1, strided, and random access

### Latency tolerance

Vectorization hides (global and local) memory latency



#### X1E weaknesses

- Limited memory per MSP
- Very slow scalar processor
  - 565 MHz
  - 2-way superscalar
  - Simple design (compared to Opteron)
- Tuning is required



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# Using Phoenix

- <a href="http://nccs.gov">http://nccs.gov</a>
  - → Resources
    - → Cray X1E Phoenix



# Login to Robin1

- 4-processor Opteron system with 32 GB of memory
  - "Robin" will point to Robin1 in a week or so
- Cross compile for Phoenix
  - Up to 30x faster than on Phoenix
- Phoenix "/tmp/work" mounted over NFS
- Most Phoenix "man" pages
- Submit and monitor jobs
  - All PBS and Moab commands
  - Even "qsub -I"; shell runs remotely on Phoenix
  - Moab command are slow; Phoenix is the server
- Typical Linux editors and tools
  - Emacs, Subversion, etc.
  - See "module avail"
- Software auto-configuration can be tricky



# Compiling

- Avoid "#ifdef CRAY"
  - Cray X1E too different from past Cray's (more like other vendors')
  - Default type sizes are not all 64 bits
- Use default optimization
  - Don't try to fix performance problems with higher optimization
- Always generate loopmarks ("-rm", "-h list=a")
- Often generate instrumented executables ("pat build")
- Try newer (or older) compilers with "module swap"
  - module avail PrgEnv
  - module swap PrgEnv PrgEnv.5509
- "-0/h gen\_private\_callee" to generate procedure interfaces for calling within CSD streams
- "-Z" for Co-Array Fortran ("-h upc" for UPC)
- "-0/h command" for serial tools on Phoenix OS nodes
  - Make sure "configure" uses this (when using Phoenix directly)



# Batch jobs

- Scheduling policy not changing (unlike Jaguar)
- Always specify requirements in MSPs
  - "-1 mppe=N"
  - For SSP jobs, divide SSP count by four
  - Jobs using more than one node (4 MSPs) must request an integer multiple of 8 MSPs
    - To line up on module boundaries for remote address translation

#### For more memory

- Tell batch system using MSP request (not memory request)
- Memory/(1.7 GB) = number of MSPs to request
- Tell "aprun" memory requirement using "-m"
- May need to set environment variables
- See "man 7 memory" on Phoenix



# Debugging

- Avoid "-g"
  - Horrible performance
  - Bugs often go into hiding
  - All levels of "-G" affect optimization
- Always set "TRACEBK" environment variable
  - seteny TRACEBK 30
  - export TRACEBK=30
- Turn on core files: "aprun -c core=unlimited"
  - Only in "/tmp/work"!
- View core files with "totalview" or "totalviewcli"
- Check traceback for hints on where to look
  - This says to look at core file #299:
    Traceback for process 64311(ssp mode) apid 64184.229 on node 7
- See online docs to try interactive debugging



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# Tuning is required

- Tuning priorities
- Vectorization
- Multistreaming
- Communication
- OpenMP?
- Tuning strategy



# Tuning priorities

- Vectorization (10x)
- Multistreaming (4x)
- Low-latency communication (2x)
- Register blocking (<2x)</li>
- Cache blocking?

#### Vectorization

- One vector instruction = many loop iterations
- Needs enough loop iterations
  - 64 (multistreamed) or 256 on X1E
  - Fewer iterations = lower efficiency
- No procedure calls
- No loop-carried data dependencies
  - Some exceptions (reductions)



# Vectorization: What the compiler can do

- Array notation
- Scalar temporary variables
- Re-arrange loop nests
- Reductions, (un)pack, scatter/gather
- Fuse loops and array statements
- Inline procedures (one level down)
- "if" statements within loops
  - Vector masks, some loss of efficiency



# Vectorization: What compilers can't do

- Make short vector loops efficient
- Make stride-1 (or -0) scatter/gather efficient
- Know that index arrays don't repeat

$$-do j = 1, n$$
  
  $x(i(j)) = x(i(j)) + ...$ 

Effectively inline many levels down

# Vectorization: How you can help

- Assert that a loop is concurrent (index arrays don't repeat)
  - -!dir\$ concurrent
  - #pragma CRI concurrent
- Assert that an index array is a permutation
  - !dir\$ permutation(i)
- Change array temporaries to scalar
  - Can remove dependencies
- Break up the big outer loop
  - To move it inside multiple inner loops
- Move loops inside procedure calls
- Move I/O outside of compute loops



# Vectorization: Loopmark listings

What vectorized, what didn't, and why?

### Beware of partial vectorization

```
6.
    Vp---- < DO i = 1,n
                e(ix1(i)) = e(ix1(i)) - a(i)
7. VP r - <>
8. VP---> END DO
f90-6371 f90: VECTOR File = gs-2.f, Line = 6
 A vectorized loop contains potential conflicts due to indirect
 addressing at line 7, causing less efficient code to be
 generated.
f90-6204 f90: VECTOR File = gs-2.f, Line = 6
 A loop starting at line 6 was vectorized.
```

#### Fix with directives

f90-6203 f90: VECTOR File = gs-2.f, Line = 7

A loop starting at line 7 was vectorized because an IVDEP or CONCURRENT compiler directive was specified.

f90-6203 f90: STREAM File = gs-2.f, Line = 7

A loop starting at line 7 was streamed because an IVDEP or CONCURRENT compiler directive was specified.

Declaring ix1 as a **permutation** may be even better



# Multistreaming

#### Compiler can multistream:

- Most vectorizable loops
- Most array syntax
- Nested loops with no dependencies
- Loop nests for vectorization within multistreaming
- Short loops

#### Compiler can't:

- Multistream loops with:
  - Procedure calls
  - Dependencies
- Always choose the right loop to vectorize versus multistream



# Multistreaming: How you can help

#### Directives, directives, directives

- -!dir\$ concurrent
- -!dir\$ preferstream
- -!dir\$ prefervector
- -!dir\$ ssp\_private
   (procedure calls)

### Cray Streaming Directives (CSDs)

- Much like OpenMP

### I/O inside a loop

```
6. 1--< do i = 1, nx
7. 1
           c(i) = a(i) * b(i)
8.
   1 write(8,'(1x,f12.4)') c(i)
9. 1--> end do
ftn-6286 ftn: VECTOR File = io1.ftn, Line = 6
 A loop starting at line 6 was not vectorized because it contains
  input/output operations at line 8.
ftn-6709 ftn: STREAM File = io1.ftn, Line = 6
 A loop starting at line 6 was not multi-streamed because it contains
  input/output operations.
```



#### Fixed

```
7. MVr--< do i = 1, nx
        c(i) = a(i) * b(i)
8.
    MVr
9. MVr--> end do
10.
11.
           write(8,'(1x,f12.4)') (c(i),i=1,nx)
ftn-6005 ftn: SCALAR File = io2.ftn, Line = 7
 A loop starting at line 7 was unrolled 2 times.
ftn-6204 ftn: VECTOR File = io2.ftn, Line = 7
 A loop starting at line 7 was vectorized.
ftn-6601 ftn: STREAM File = io2.ftn, Line = 7
 A loop starting at line 7 was multi-streamed.
```

#### Communication

- Use one-sided communication for latency-sensitive operations
- MPI-2 library
  - Complicated interface
  - No guaranteed progress without synchronization
- SHMEM library
  - Vendor specific
- Co-Array Fortran
  - Lowest latency
  - Currently vendor specific
  - Part of next Fortran language standard
- Intermix with each other and MPI-1



# OpenMP?

### If OpenMP used for different parallelism than MPI

- Probably the same parallelism as for vectorization and multistreaming
- Typically not enough parallel work for all three
- OpenMP is least efficient of the three

### If OpenMP used for same parallelism as MPI

- Useful for reducing message volume and aggregating messages
- But one MSP can't saturate the network
- Little reason to aggregate
- Don't bother with OpenMP on Phoenix



# Tuning strategy

- Functional port
- Iterate
  - Loopmark and profile
  - Vectorize and multistream
- Tune communication



# Profiling

- Instrument executable with "pat\_build"
- Run and generate performance report
  - Together: "pat\_run"
  - Separately: "aprun" followed by offline "pat\_report" on resulting ".xf" file
- Use report to locate bottlenecks, then use loopmark listings to diagnose problems and solutions
  - Use call-tree reports to find which calls were expensive
  - Apprentice2 tool provides graphical browsing
- For aggregate hardware-counter statistics, use "pat\_hwpc"



#### More information

- This information and more at:
   <u>http://nccs.gov</u> → Resources → Cray X1E Phoenix
   <u>http://info.nccs.gov</u> → Cray X1E Phoenix
   <u>http://info.nccs.gov/resources/phoenix</u>
- Cray documentation at <a href="http://docs.cray.com">http://docs.cray.com</a>
  - Cray X1 Series System Overview
  - Cray Fortran, C/C++ reference manuals
  - Migrating Applications to the Cray X1 Series Systems
  - Optimizing Applications on Cray Series Systems
  - Excellent search capability
- E-mail us: <u>help@nccs.gov</u>

